



seL4[®] verification: status and plans

Michael McInerney @ Proofcraft



Key-se14



The world's most highly assured operating system kernel



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Unparalleled mathematical proofs
of correctness and security

⇒ The most trustworthy foundation
for safety- and security-critical systems



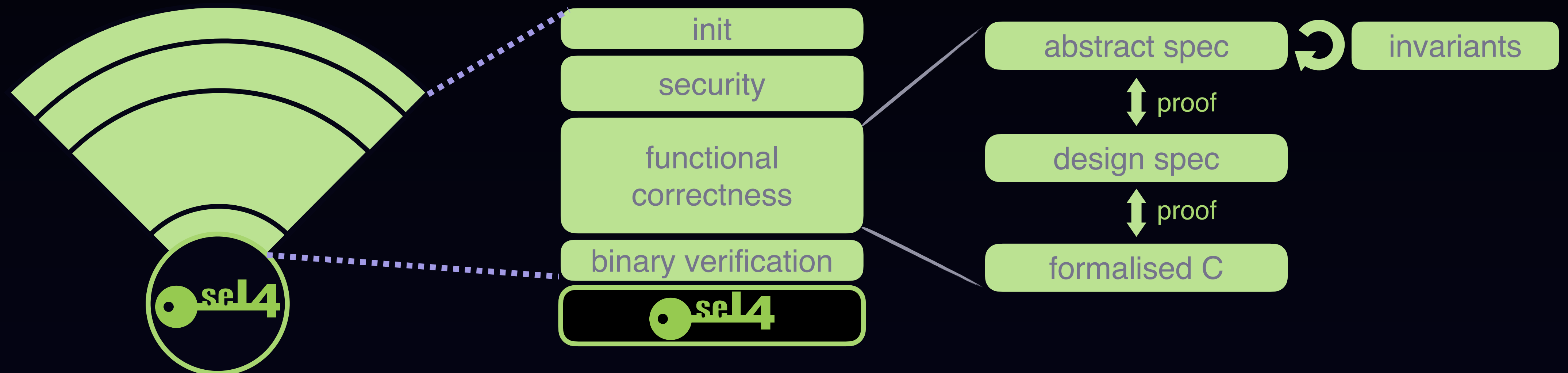
Unparalleled mathematical proofs
of correctness and security

seL4

Unparalleled mathematical proofs
of correctness and security

(FC)

(integrity&confidentiality)





Unparalleled mathematical proofs
of correctness and security

NOW & IN THE FUTURE

More architectures verified

More features verified

More platforms verified

More cores

Less need for
expertise & maintenance

Overview: 5 main areas Proofcraft is working on



More architectures verified

Arm 64-bit (AArch64)

More features verified

Mixed-criticality support (MCS)

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Automated platform port verification

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Proof architecture split (arch-split)

More cores

Verified multikernel (MK) on multicore

Overview: status & plans in a nutshell



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Now: 18/32 configs verified (56%)
Aim: 90-100% for existing ones
+ automation for new ones

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seL4 proofs

Done
Ongoing
Future

(non-MCS, uncore)

Done: FC
Now: integrity (Q1'25)

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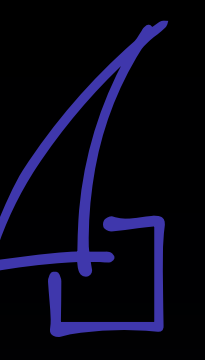
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What is MCS?



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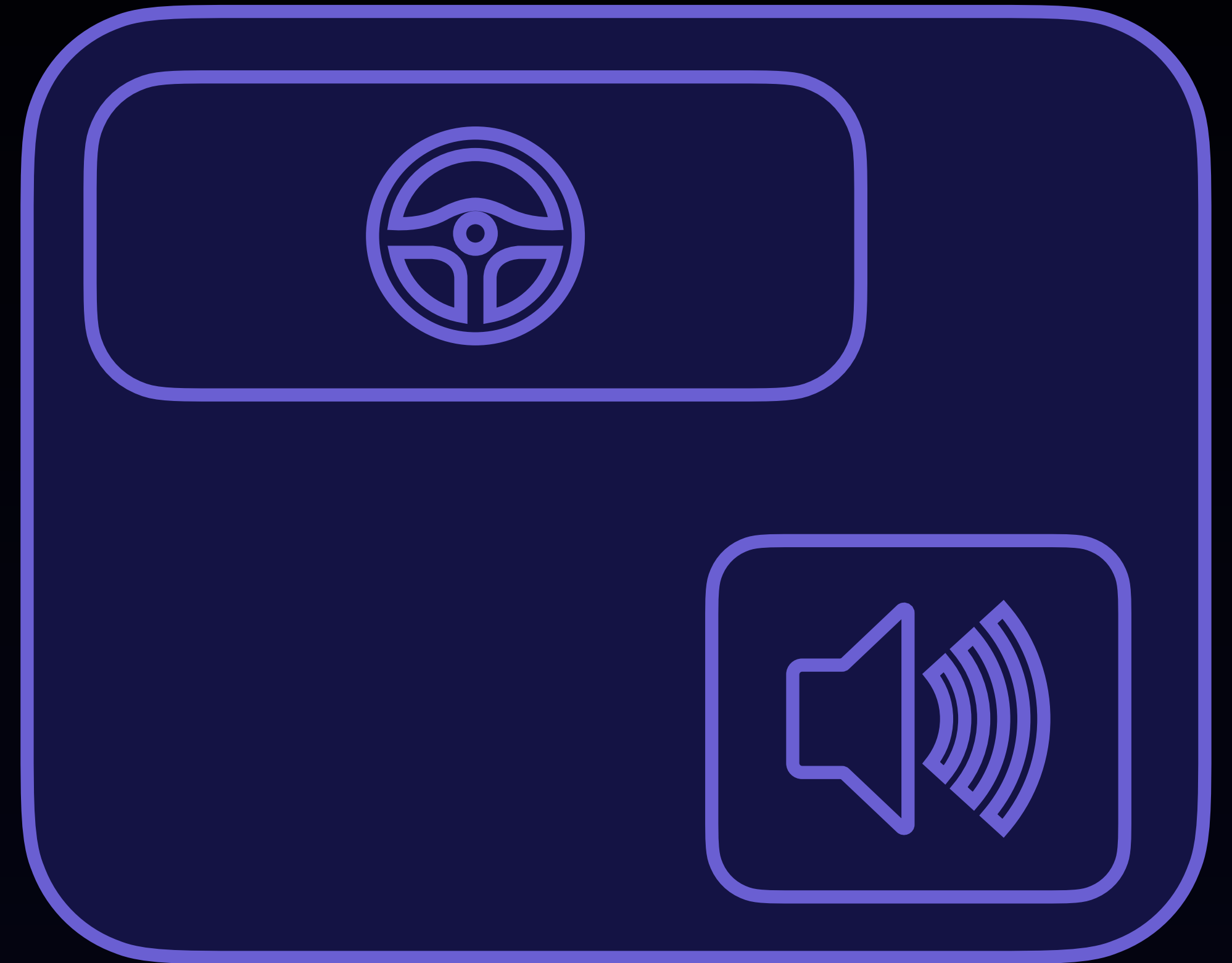


- Support for Mixed-Criticality Systems

What is MCS?



- Support for Mixed-Criticality Systems



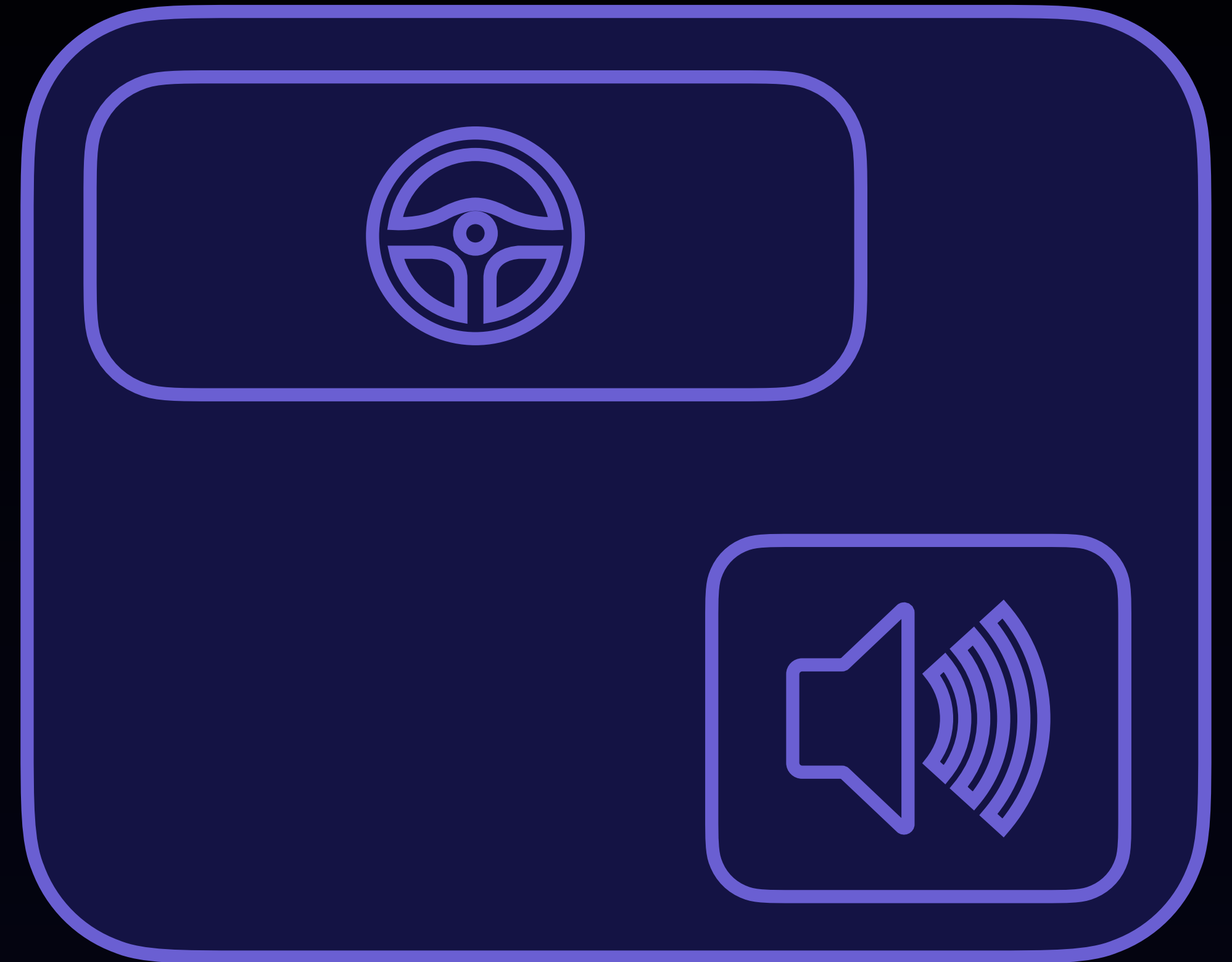
What is MCS?

- Support for Mixed-Criticality Systems
- Time as a resource



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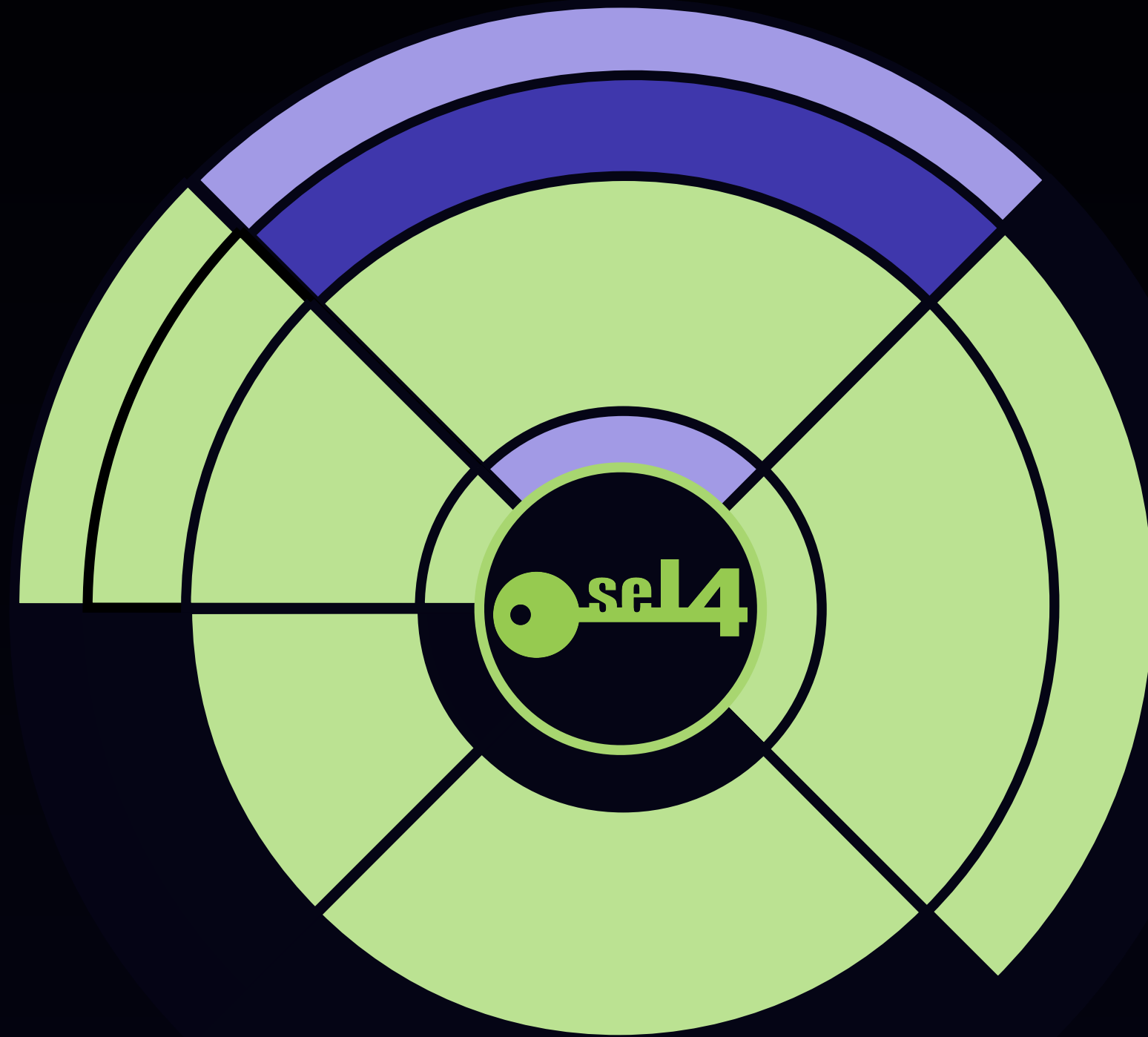
- Support for Mixed-Criticality Systems
- Time as a resource
 - scheduling context objects



The proofs have evolved with new features over the years



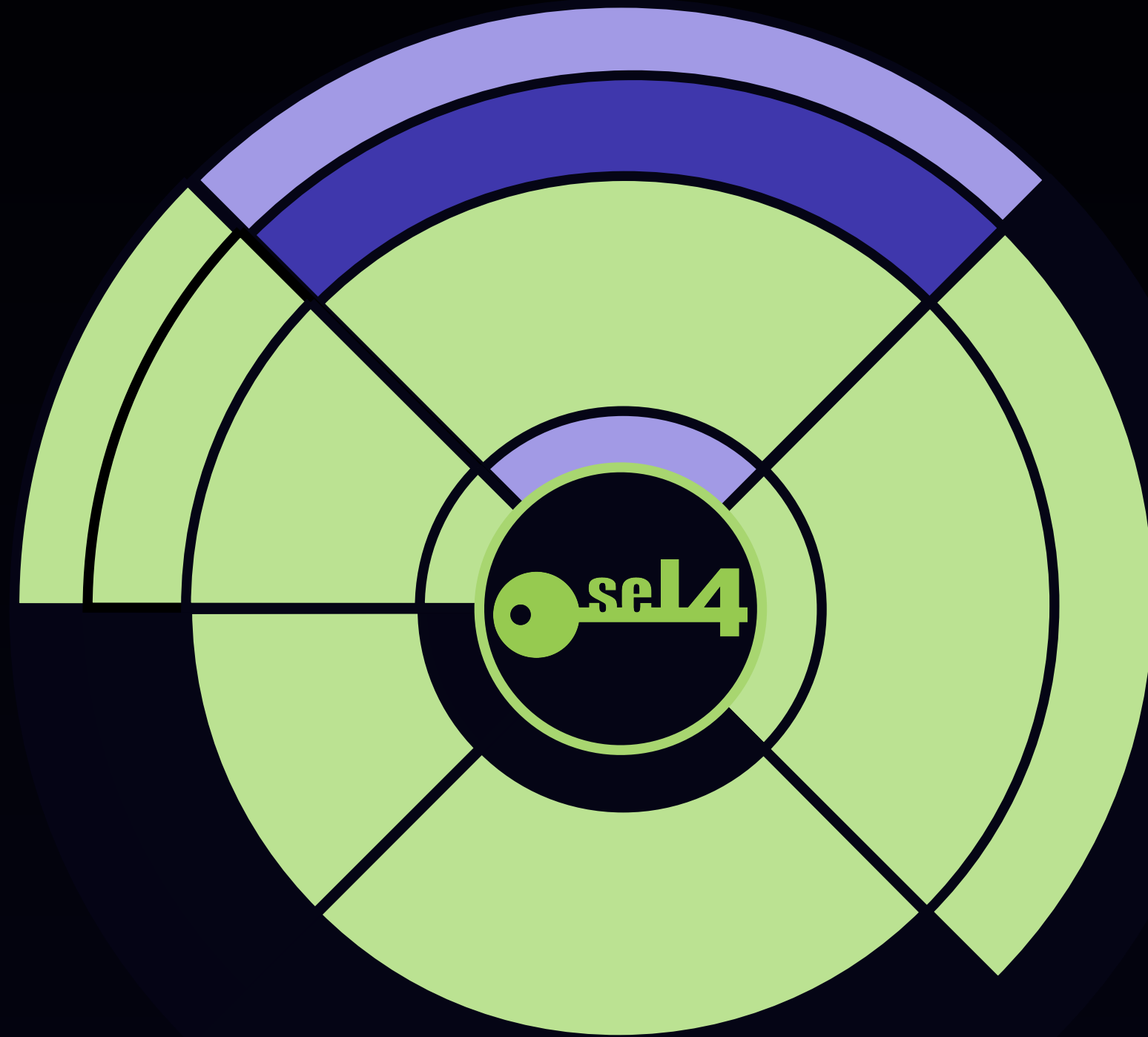
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Two examples:

- bound notification endpoints
- bitfield scheduler optimisation

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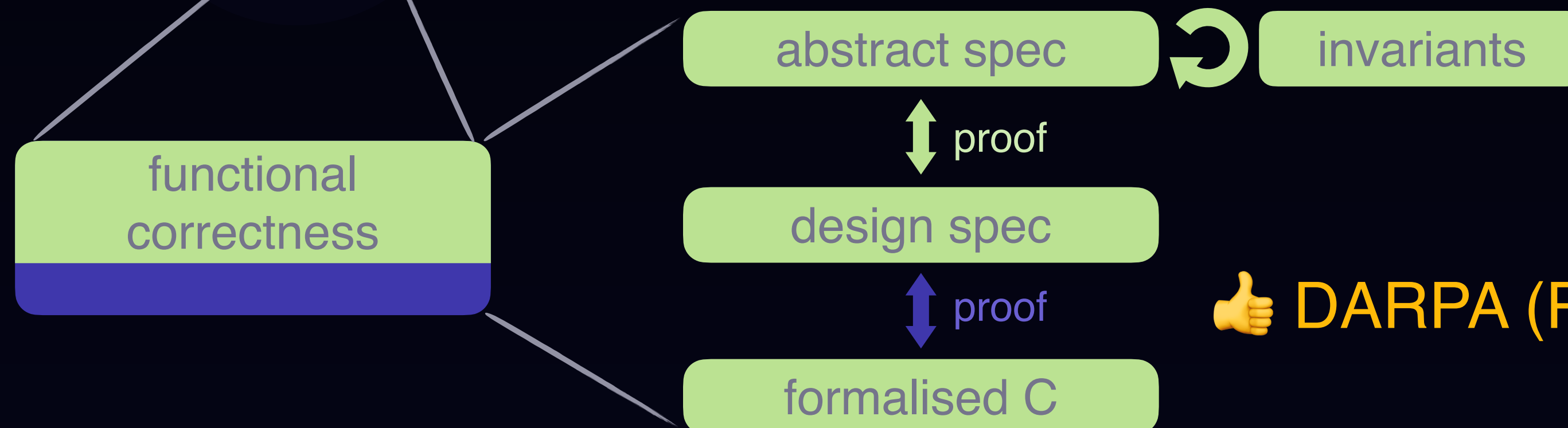
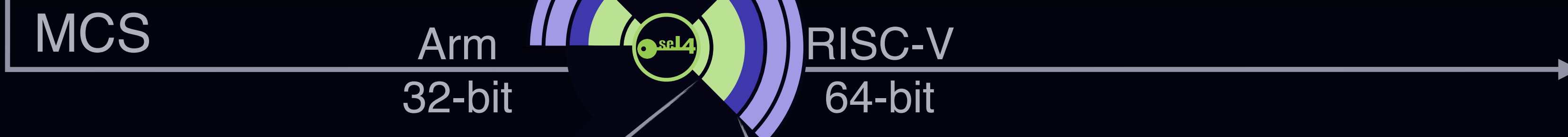
MCS is different:

- large, invasive change

Big Feature: Mixed-Criticality Systems



Verification of multiple configs in parallel



👍 DARPA (PROVERS)

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- Many pre-existing functions are now much longer

Loops in MCS



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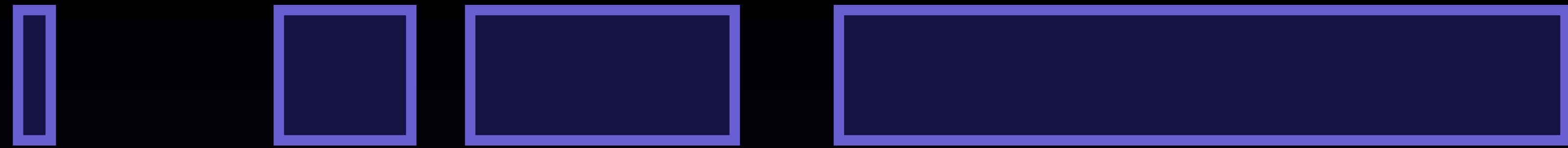
Loops in MCS



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Loops in MCS



t_0

t_1

t_2

t_3



Time

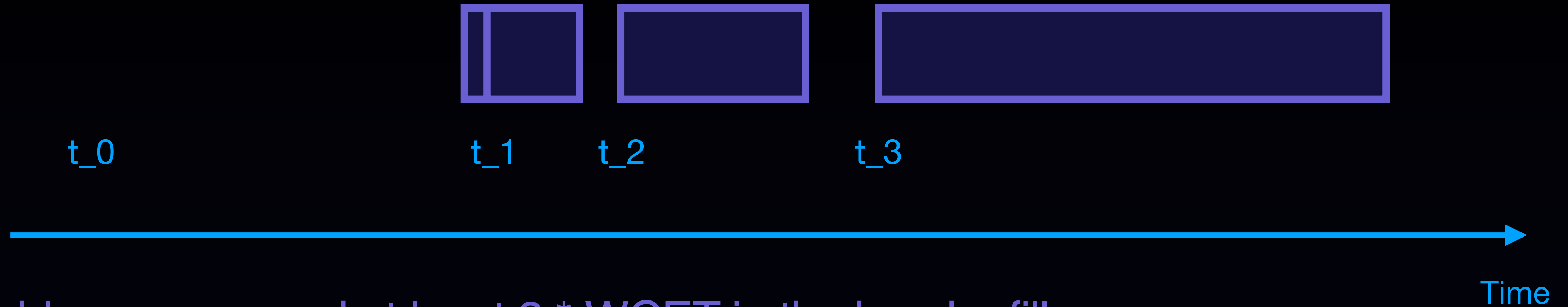
Problem...

Loops in MCS



Problem... we need at least $2 * WCET$ in the head refill

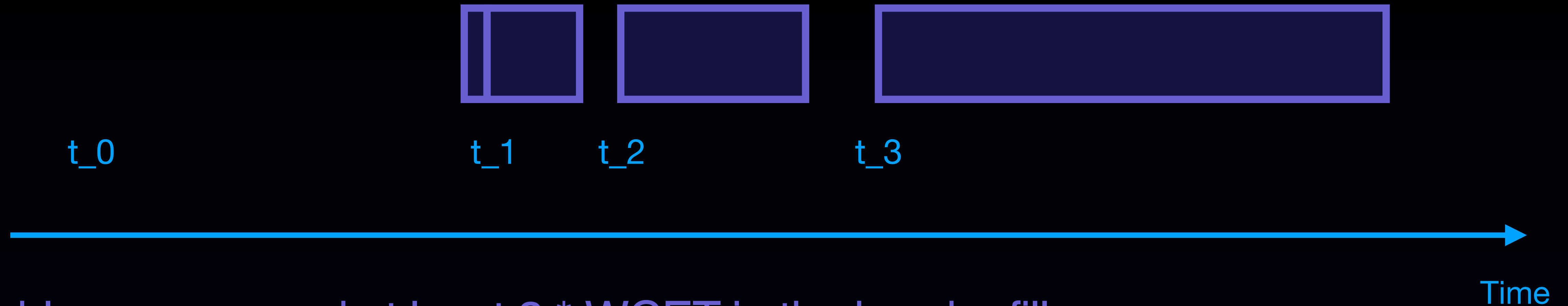
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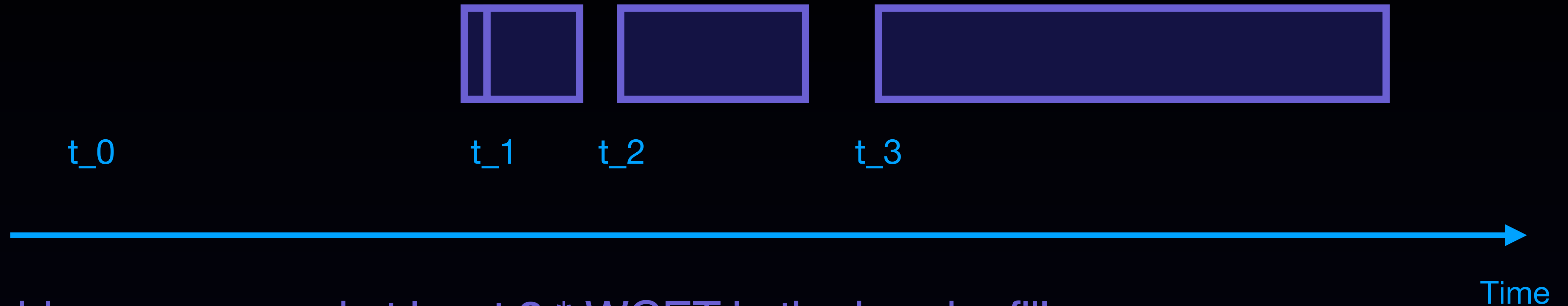
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Loops in MCS

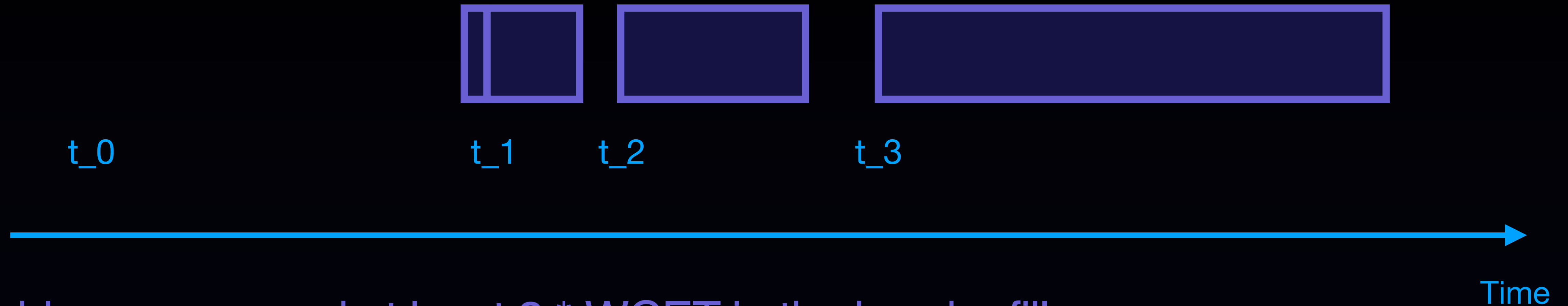


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Data refinement? Ring buffer of refills versus list of refills

Conclusion



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